

Formal Languages applied to Linguistics

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Overview

- 1 Formal Languages
- 2 Formal Grammars
- 3 Regular Languages
 - Definition
 - Automata
 - Properties
- 4 Formal complexity of Natural Languages

Definition

3 possible definitions

- 1 a regular language can be generated by a regular grammar
- 2 a regular language can be defined by rational expressions
- 3 a regular language can be recognized by a finite automaton

Def. 15 (Rational Language)

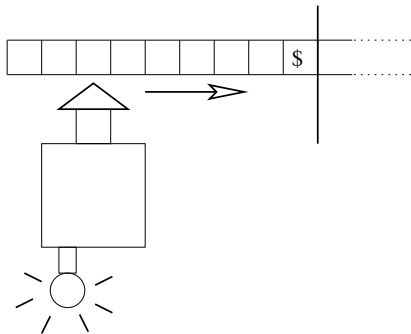
A rational language on Σ is a subset of Σ^* inductively defined thus:

- \emptyset and $\{\varepsilon\}$ are rational languages ;
- for all $a \in X$, the singleton $\{a\}$ is a rational language ;
- for all g and h rational, the sets $g \cup h$, $g.h$ and g^* are rational languages.

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Metaphoric definition



Formal definition

Def. 16 (Finite deterministic automaton (FDA))

A finite state deterministic automaton \mathcal{A} is defined by :

$$\mathcal{A} = \langle Q, \Sigma, q_0, F, \delta \rangle$$

Q is a finite set of states

Σ is an alphabet

q_0 is a distinguished state, the initial state,

F is a subset of Q , whose members are called
final/terminal states

δ is a mapping **fonction** from $Q \times \Sigma$ to Q .

Notation $\delta(q, a) = r$.

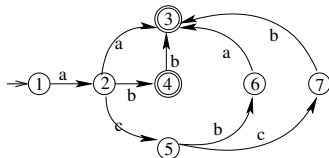
Example

Let us consider the (finite) language $\{aa, ab, abb, acba, accb\}$.

The following automaton recognizes this language: $\langle Q, \Sigma, q_0, F, \delta \rangle$,
 avec $Q = \{1, 2, 3, 4, 5, 6, 7\}$, $\Sigma = \{a, b, c\}$, $q_0 = 1$, $F = \{3, 4\}$, and
 δ is thus defined:

δ :

- (1,a) \mapsto 2
- (2,a) \mapsto 3
- (2,b) \mapsto 4
- (2,c) \mapsto 5
- (4,b) \mapsto 3
- (5,b) \mapsto 6
- (5,c) \mapsto 7
- (6,a) \mapsto 3
- (7,b) \mapsto 3



| | a | b | c |
|-----|---|---|---|
| → 1 | 2 | | |
| 2 | 3 | 4 | 5 |
| ← 3 | | | |
| ← 4 | | 3 | |
| 5 | | 6 | 7 |
| 6 | 3 | | |
| 7 | | 3 | |

Recognition

Recognition is defined as the existence of a sequence of states defined in the following way. Such a sequence is called a path in the automaton.

Def. 17 (Recognition)

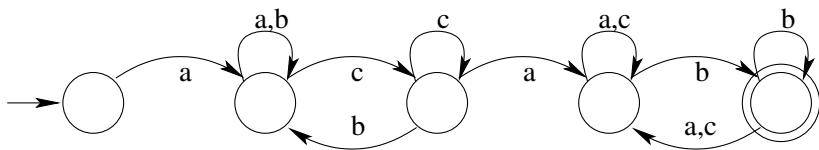
A word $a_1a_2\dots a_n$ is **recognized/accepted** by an automaton iff there exists a sequence k_0, k_1, \dots, k_n of states such that:

$$k_0 = q_0$$

$$k_n \in F$$

$$\forall i \in [1, n], \delta(k_{i-1}, a_i) = k_i$$

Example



Exercices

Let $\Sigma = \{a, b, c\}$. Give deterministic finite state automata that accept the following languages:

- 1 The set of words with an even length.
- 2 The set of words where the number of occurrences of b is divisible by 3.
- 3 The set of words ending with a b .
- 4 The set of words not ending with a b .
- 5 The set of words non empty not ending with a b .
- 6 The set of words comprising at least a b .
- 7 The set of words comprising at most a b .
- 8 The set of words comprising exactly one b .

Answers

